## Adobe Source Libraries Overview \& Philosophy

## Sean Parent

Principal Scientist \& Engineering Manager Adobe Software Technology Lab
http://stlab.adobe.com
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Demo

## Adobe Source Libraries

- A collection of libraries to support application development
- Research artifacts of the Adobe Software Technology Lab
- Open Source: http://stlab.adobe.com/
- Used by many Adobe products


## Outline

- Regular Types - libraries for efficiently handling regular types
- Forest - advantages of explicit data structures
- Layout Library - a library for placing / aligning items in an interface and a language to express layouts
- Property Model Library - describing and solving inter-related proprerties


## Goal of ASL

- Express entire applications using a combination of generic and declarative techniques
- 2 orders of magnitude reduction in code
- Greater than corresponding reduction in defects
- We are still a long way from our goal
- perhaps not as far as it would appear


## Approach

- Generic Algorithms
- Write algorithms with minimal requirements - maximum reuse
- Generic Data Structures (Containers)
- Containers support algorithm requirements (including complexity)
- Declarative Architecture
" Identify "patterns" of how components are assembled and learn to express/solve these pattern with algorithms and data structures


## Challenges

- Build a Strong Foundation
- See http://stepanovpapers.com/eop/lecture alll.pdf
- Our work here has a strong impact on all aspects of ASL
- Combine Runtime Polymorphism and Generic Programming
- See http://www.emarcus.org/papers/gpce2007f-authors-version.pdf
- See http://www.emarcus.org/papers/MPOOL2007-marcus.pdf
- See Poly and Any Regular Libraries
- Make Implicit Structure Explicit
- Work ongoing - see Forest, Property Model, and Layout Libraries
- Discovering the Rules that Govern Large Systems
- Work ongoing - see Property Model Library and initial work on Sequence Models

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## Adobe Source Libraries - Regular Types

- Definition: Regular
- Move Library
- How RVO works
- Creating Polymorphic Regular Types and Poly Library
- Copy On Write Library


## Definition of Regular

- The requirements of Regular are based on equaltional reasoning
- They assure regularity of behavior and interoperability
- Types which model these requirements are regular types
- The properties of Regular are inherent in the machine model
- Regular types exist in any correct system but formalizing the requirements and normalizing the syntax is what enables interoperability
- All types are inherently regular


## Basic Requirements of Regular Type

| Requirement | Syntax Example | Axioms \& Postconditions |
| :---: | :---: | :---: |
| Copy | $\begin{aligned} & \mathrm{T} x=y ; \\ & \sim x() ; \end{aligned}$ | $\begin{aligned} & x==y \\ & \text { if (is_defined(modify, x) } \\ & \quad \text { then modify }(x) ; x!=y \end{aligned}$ |
| Assignment | $x=y ;$ | $\begin{aligned} & x==y \\ & \text { if (is_defined(modify, x) } \\ & \quad \text { then modify }(x) ; x!=y \end{aligned}$ |
| Equality | $\begin{aligned} & x==y ; \\ & x \quad!=y \end{aligned}$ | $\begin{aligned} & a==b \& \& b==c=a=c \\ & a=b \Leftrightarrow b=a \\ & a==a \end{aligned}$ |
| Identity | \& F | ```&a == &b => a == b given &x == &y if (is_defined(modify, x) then modify(x); x == y;``` |
| Size | sizeof(T); | size of local part of $T$ |
| Swap | $\operatorname{swap}(\mathrm{x}, \mathrm{y})$; | $x^{\prime}==y ; y^{\prime}==x$; <br> O(sizeof(T)); nothrow; |

## Extended Requirements of Regular Type

| Requirement | Syntax Example | Axioms \& Postconditions |
| :---: | :---: | :---: |
| Default Construction | T x; | T $x$; $x=y$; is equivalent T $x=y$; |
| Default Comparison | $\begin{gathered} \text { std: }: \text { less }<T>() \\ (x, y) ; \end{gathered}$ | $\begin{gathered} !o p(x, y) \& \&!o p(y, x) \\ \Rightarrow x==y \end{gathered}$ |
| Movable | $\begin{aligned} & x=f() ; \\ & x=\operatorname{move}(y) \end{aligned}$ | $\begin{aligned} & \text { O(sizeof }(\mathrm{T})) \text {; nothrow; } \\ & \mathrm{T} \mathrm{x}=\mathrm{y} ; \mathrm{z}=\operatorname{move}(\mathrm{x}) \\ & =>\mathrm{z}==\mathrm{y} \end{aligned}$ |
| Area | $\operatorname{area}(\mathrm{x})$; | Copy and Assignment are $0(\operatorname{area}(x))$; <br> Equality is worst case $0(\operatorname{area}(x))$; |
| Alignment | alignment(T); | alignment size for type |
| Underlying Type | underlying(T) | type which can be copied to/ from $T$ in $0(\operatorname{size}(T))$ |

## Importance of Move

- Allows transfer of ownership of remote parts in small constant time
- Will not throw an exception
- Move does not refine Copy and Copy does not refine Move
- When the source will not be used after a copy, copy can be replaced with move
- An object which has been moved from is still Regular
- Reference Semantics provide move for "free"
- But there are other costs

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## Quiz: What will the following code print?

```
struct object_t
{
    object_t()
        { cout << "construct" << endl; }
    object_t(const object_t&)
        { cout << "copy" << endl; }
    object_t& operator=(const object_t&)
        { cout << "assign" << endl; return *this; }
};
object_t function()
    { object_t result; return result; }
int main()
    { object_t x = function(); return 0; }
```


## Answer: Return Value Optimization Eliminates Copies

```
struct object_t
{
    object_t()
        { cout << "construct" << endl; }
    object_t(const object_t&)
        { cout << "copy" << endl; }
    object_t& operator=(const object_t&)
        { cout << "assign" << endl; return *this; }
};
object_t function()
    { object_t result; return result; }
int main()
    { object_t x = function(); return 0; }
```construct

\section*{Quiz: What will the following code print?}
```

struct object_t
{
object_t()
{ cout << "construct" << endl; }
object_t(const object_t\&)
{ cout << "copy" << endl; }
object_t\& operator=(const object_t\&)
{ cout << "assign" << endl; return *this; }
};
object_t function()
{ object_t result; return result; }
void sink(object_t) { }
int main()
{ sink(function()); return 0; }

```

\section*{Answer: RVO Works for Parameters Also}
```

struct object_t
{
object_t()
{ cout << "construct" << endl; }
object_t(const object_t\&)
{ cout << "copy" << endl; }
object_t\& operator=(const object_t\&)
{ cout << "assign" << endl; return *this; }
};
object_t function()
{ object_t result; return result; }
void sink(object_t) { }
int main()
{ sink(function()); return 0; }

```
construct

\section*{Sink Functions}
- A sink function is any function which consumes one or more arguments by storing them or by returning them
- By passing the argument by value and moving it into position we allow the compiler to avoid a copy
- Assignment is a sink function

\section*{Typical Assignment}
```

struct object_t{
object_t() : object_m(new int(0)) { }
object_t(const object_t\& x) : object_m(new int(*x.object_m))
{ cout << "copy" << endl; }
object_t\& operator=(const object_t\& x)
{ object_t tmp = x; swap(tmp, *this); return *this; }
~object_t() { delete object_m; }
friend inline void swap(object_t\& x, object_t\& y)
{ swap(x.object_m, y.object_m); }
private:
int* object_m;
};
object_t function()
{ object_t result; return result; }
int main()
{ object_t x; x = function(); return 0; }
copy

```

\section*{Better Assignment}
```

struct object_t{
object_t() : object_m(new int(0)) { }
object_t(const object_t\& x) : object_m(new int(*x.object_m))
{ cout << "copy" << endl; }
object_t\& operator=(object_t x)
{ swap(x, *this); return *this; }
~object_t() { delete object_m; }
friend inline void swap(object_t\& x, object_t\& y)
{ swap(x.object_m, y.object_m); }
private:
int* object_m;
};
object_t function()
{ object_t result; return result; }
int main()
{ object_t x; x = function(); return 0; }

```
copy

\section*{Better Assignment}
```

struct object_t{
object_t() : object_m(new int(0)) { }
object_t(const object_t\& x) : object_m(new int(*x.object_m))
{ cout << "copy" << endl; }
object_t\& operator=(object_t x)
{ swap(x, *this); return *this; }
~object_t() { delete object_m; }
friend inline void swap(object_t\& x, object_t\& y)
{ swap(x.object_m, y.object_m); }
private:
int* object_m;
};
object_t function()
{ object_t result; return result; }
int main()
{ object_t x; x = function(); return 0; }

```

\section*{Explicit Move}
```

struct object_t{
object_t(move_from<object_t> x) : object_m(0)
{ swap(*this, x.source); }
int\& get() { return *object_m; }
//..
};
object_t function()
{ object_t result; return result; }
object_t sink(object_t x)
{ x.get() += 5; return move(x); }
int main()
{ object_t x = sink(function()); return 0; }

```

\section*{Polymorphism and Regular Types}
- Current pattern:
- polymorphism => inheritance => specialized classes => limited code sharing
- polymorphism => variable size => heap allocation => pointer management
- polymorphism => virtual functions => slower dispatch
- The requirement for polymorphism comes from the need to handle heterogeneous types which satisfy a common set of requirement in a homogeneous manner
- Requirement is driven by the use of the type, there is nothing inherently polymorphic about a type

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\section*{Creating a Polymorphic Regular Type}
```

struct object_t
{
template <typename T> // T models Drawable
explicit object_t(T x) : object_m(new model_t<T>(move(x))) { }
object_t(move_from<object_t> x) : object_m(0)
{ swap(*this, x.source); }
object_t(const object_t\& x) : object_m(x.object_m->copy_()) { }
object_t\& operator=(object_t x) { swap(x, *this); return *this; }
~object_t() { delete object_m; }
friend inline void swap(object_t\& x, object_t\& y)
{ using std::swap; swap(x.object_m, y.object_m); }
friend inline void draw(const object_t\& x)
{ x.object_m->draw_(); }
private:
// ...fill in here...
concept_t* object_m;
};

```

\section*{Creating a Polymorphic Regular Type}
```

struct concept_t
{
virtual ~concept_t() { }
virtual concept_t* copy_() const = 0;
virtual void draw_() const = 0;
};
template <typename T>
struct model_t : concept_t
{
explicit model_t(T x) : value_m(move(x)) { }
concept_t* copy_() const { return new model_t(*this); }
void draw_() const { draw(value_m); }
T value_m;
};

```

\section*{Using our Poly Drawable Type}
```

template <typename T> void draw(const T\& x) { cout << x << endl; }
template <typename T> void draw(const vector<T>\& x) {
typedef typename vector<T>::const_iterator iterator_t;
cout << "<vector>" << endl;
for (iterator_t f(x.begin()), l(x.end()); f != l; ++f)
{ draw(*f); }
cout << "</vector>" << endl;
}
int main() {
vector<object_t> x;
x.push_back(object_t(10));
x.push_back(object_t(string_t("Hello World!")));
x.push_back(object_t(x));
x.push_back(object_t(string_t("Another String!")));
draw(x);
return 0;
}

```

\section*{Results}
```

<vector>
    10
    Hello World!
    <vector>
        10
        Hello World!
    </vector>
    Another String!
</vector>
```

\section*{Indenting Added for clarity}

\section*{Summary}
- Non-Intrusive - client need only satisfy requirements
- Existing types can be used in a polymorphic fashion without wrapping
- Cost of virtual dispatch the same - but only required when object used in a polymorphic setting
- Client isn't burdened by managing pointers - can use efficiently with containers and algorithms
- The Poly Library provides facilities for:
- Virtualization of the properties of Regular
- Refinement
- Dynamic Type Information

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\section*{One Final Change...}
```

template <typename T>
void draw(const copy_on_write<T>\& x) { draw(x.read()); }
int main(){
typedef copy_on_write<object_t> cow_t;
vector<cow_t> x;
x.push_back(cow_t(object_t(10)));
x.push_back(cow_t(object_t(string_t("Hello World!"))));
x.push_back(cow_t(object_t(x)));
x.push_back(cow_t(object_t(string_t("Another String!"))));
draw(x);
return 0;
}

```

\section*{Forest Library}
- STL provides sequence and associative containers and algorithms
- Because the STL data types are Regular they can be composed to create new structures
- Not all structures are best represented by composition
- Hierarchies can be represented through containment
- as we saw with object_t
- Other representations provide other advantages

\section*{Forest}


\section*{Forest (full-order traversal)}


\section*{Forest (pre-order traversal)}


\section*{Forest (post-order traversal)}


\section*{Forest (child traversal)}

\[
\begin{array}{lll}
\phi 0 \phi & \$ & \phi \\
0 & 00
\end{array}
\]

\section*{Print as XML}
```

template <typename T> // T models Regular
ostream\& operator<<(ostream\& stream, const forest<T>\& x)
{
typedef typename forest<T>::const_iterator iterator_t;
typedef depth_fullorder_iterator<iterator_t> depth_iterator_t;
for (depth_iterator_t f(begin(x)), l(end(x)); f != l; ++f)
{
for (size_t n(f.depth()); n != 0; --n) stream << "\t";
stream << (f.edge() ? "<" : "</") << *f << ">" << endl;
}
return stream;
}

```

\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << X;
return 0;
}

```

\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << x;
return 0;
}

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << x;
return 0;
}

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
D
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << x;
return 0;
}

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << x;
return 0;
}

```
```

    ++i;
    ```
    ++i;
iterator_t j = x.insert(i, "son");
```

iterator_t j = x.insert(i, "son");

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << X;
return 0;
}

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << X;
return 0;
}

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i,, "daughter");
cout << X;
return 0;
}

```


\section*{Example}
```

int main()
{
typedef forest<const char*> forest_t;
typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++j;
x.insert(j, "grandson");
x.insert(i, "daughter");
cout << X;
return 0;
}

```


\section*{Example}
int main()
\{
typedef forest<const char*> forest_t; typedef forest_t::iterator iterator_t;
forest_t x;
iterator_t i = x.insert(x.end(), "me");
x.insert(x.end(), "brother");
++i;
iterator_t j = x.insert(i, "son");
++ ;
x.insert(j, "grandson");
x.insert(i, "daughter");cout << x;
return 0;
\}
```

<me>
    <son>
                <grandson>
                </grandson>
    </son>
    <daughter>
    </daughter>
</me>
<brother>
</brother>
```


\section*{Declarative Ul with ASL}
- Introduction
- What a User Interface Is
- Identifying UI Mechanisms
- What MVC Is
- Property Models and Layouts Libraries
- Modeling the Form
- Presenting the Form
- Property Model Basics
- An Overview of The Property Model Syntax
- CEL expression and the Begin Inspector
- Invariants \& Dependency Tracking
- Relationships \& Logic
- Layout Library Basics
- An Overview of the Layout Library Syntax
- Placement and Alignment
- Spacing, Margins, and Indenting
- Guides
- Optional and Panel
- Advanced Topics
- Scripting and Localization
- How Layouts Work
- What you can't do
- How Property Models Work
- What you can't do

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\section*{What is a User Interface?}

\section*{Discussion}

\section*{What a User Interface Is}
- Definition: A User Interface (UI) is a system for assisting a user in selecting a function and providing a valid set of parameters to the function.
- Definition: A Graphical User Interface (GUI) is a visual and interactive UI.

\section*{Mechanisms to Assist the User}

\section*{Discussion}

\section*{UI Mechanisms}
- Context
- Current Document, Selection, Tools, Modal Dialogs
- Context Provides a Function or One or More parameters to the Function
- The current item is referred to as the subject
- The selected function is the verb
- Sentences
- subject-verb(function)-[object]
- Drag and Drop, Cut/Copy/Paste
- Constraints
- Disabled Options, Rejecting Invalid Input, Modality
- Consistency

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\section*{UI Mechanisms (Continued)}
- Interactivity
- Tracking: \(\quad \approx 1 / 30 \mathrm{~s}\)
- Acknowledge: \(\approx \mathbf{1 / 5}\) s
- Confirmation: \(\approx 1 \mathrm{~s}\)
- Precognition
- Specifying Parameters in Terms of Desired Results:
- Compress this movie to fit on a DVD
- Scale this image to fit the Page
- Time-Travel
- Undo, Preview, Non-Destructive Editing
- Metaphors
- Using knowledge transference

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\section*{Introduction}

\section*{Demo}

\section*{Model-View-Controller}
- View \& Controller Logically Separate
- Most Descriptions get MVC Wrong - see Design Patterns or Smalltalk, not Apple or Microsoft.
- CMV Would be a Better Term


\section*{Model View Controller}


\section*{Model-View-Controller}


\section*{Property Models and Layouts Libraries}
- Property Model Library is only concerned with the model portion
- It is not the only way to construct a model
- Layout Library is only concerned with how the view portions are positioned in a coordinate space
- Within our Layout Descriptions we'll also providing binding to connect the widgets to the model
- It is important to note that the layout library does not have any built in knowledge about the widgets - we provide a sample set of widgets but they are not complete implementations.

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\section*{Relation to MVC}


\section*{Property Model Basics}

\section*{Property Model Descriptions}
```

sheet my_sheet
{
interface:
team_1: "Giants";
team_2: "Patriots";
score_1: 0;
score_2: 0;
output:
result <== {
team_1: team_1, team_2: team_2,
score_1: score_1, score_2: score_2 };
}

```

\section*{Property Model Descriptions}
- Interface Cells
- Optional Initializer and Expression
score_1: \(0<==\) score_2 * 2;
- Output Cells
- Require Expression
result <== [score_1, score_2];

\section*{CEL Expressions}
- Built-In Data Types
- number: -17.3
- string:
"Hello" ' world!'
- name:
@identifier
- boolean:
true
- array:
[false, "Test", @key]
- dictionary:
- empty:
\{key_1: "Value", key_2: 10\}
empty
- Variables and Function
- variable: score_1
- function: \(\max (10\), score_1)
scale(m: base, x: 10, b: offset)

\section*{CEL Expressions}
- Operators
- number: \(\quad *, /,+,-\)
- number: \(<,>,<=,>=\)
- boolean: !, \&\&, ||
- any:
==, ! =
- array:
[number_expression]
- dictionary:
- any:
[name_expression], .
expression ? expression : expression
- empty:
empty
- C order of Precedence
- Example
\{ width: 10, height: 20 \}[ p ? @width : @height]

\section*{Property Model Descriptions}
- Invariant Cells
- Requires Boolean Expression
invariant:
check <== \(a<b\);
- The pre-conditions to a function are an invariant of the functions arguments
- Cells that contribute to an invariant are poison
- Cells derived from poison are invalid

\section*{Property Model Descriptions}
- Logic Cells
- Requires Expression
logic:
\[
\text { rate }<==a^{*} b \text {; }
\]
- A logic cell is simply a named expression
- Relate Expression
logic:
```

relate {
a <== b * c;
b <== a / c;
c<== a / b;
}

```
- N-Way, Exactly One Expression Is Executed For A Given State

\section*{Visualizing Property Models}


\section*{Mini-Image Size Example}


\section*{Declarative Solution using the Property Model Library}
```

sheet mini_image_size
{
input:
original_width : 5 * 300;
original_height : 7 * 300;
interface:
constrain : true;
width_pixels : original_width <== round(width_pixels);
height_pixels : original_height <== round(height_pixels);
width_percent;
height_percent;
logic:
relate {
width_pixels <== round(width_percent * original_width / 100);
width_percent <== width_pixels * 100 / original_width;
}
relate {
height_pixels <== round(height_percent * original_height / 100);
height_percent <== height_pixels * 100 / original_height;
}
when (constrain) relate {
width_percent <== height_percent;
height_percent <== width_percent;
}
output:
result <== { height: height_pixels, width: width_pixels };
}

```

\section*{Imperative Solution to Mini-Image Size}


\section*{Event Flow in a Simple User Interface}


\section*{Layout Library Basics}

\section*{Layout Description}
```

layout my_dialog
{
interface:
display : true;
constant:
dialog_name : "My Dialog";
view dialog(name: dialog_name) {
reveal(name: "Display", bind: @display);
optional(bind: @display) {
button(name: "OK");
}
}
}

```

\section*{Placement and Alignment}
- Placement is a container property
- placement: place_row, place_column, place_overlay
- The containers row(), column(), and overlay() are non-creating containers with the corresponding placement.
- Alignment is a general property that applies to horizontal and vertical
- horizontal: align_left, align_right, align_center, align_proportional, align_fill
- vertical: align_top, align_bottom, align_center, align_proportional, align_fill
- Alignment of children can be imposed from container
- child_horizontal:
- child_vertical:
- Tip: If widgets are stuck top/left, it is likely because the container they are in isn't using align_fill.

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\section*{Spacing, Margins, Indenting}
- Spacing is a container property
- spacing: number
- spacing: array
- The spacing between each element in the container
- Margin is a container property
- margin: number
- margin: [top, left, bottom, right]
- Indent is a general property
- Indent: number
- The indent applies to the horizontal position of an item in a column and vertical position of an item in a row and is relative to the left or right alignment
- Tip: Define meaningful constants for these elements - don't use raw values and don't use to "fake" alignment.

\section*{Guides}
- Guides are Defined By Widgets (Currently)
- There are (Currently) Two Guide Types: @guide_baseline, @guide_label
- Guides Propagation Can Be Suppressed:
- guide_mask: [@guide_xxxx]
- The default mask for columns is [@guide_baseline]
- Guides Can Also Be Balanced Within A Container
- guide_balance: [@guide_xxxx]
- Guides only apply to items which are aligned left/right or top/bottom or filled. Fill left or right is determined by widget (and may vary by local).
- Tip: Guides can be allowed to propagate from overlays to get consistent column widths on tab panels.

\section*{Optional and Panel}
- optional() and panel() are containers whose visibility can be bound
- An optional() container is removed from the layout when hidden
- A panel() remains part of the layout when hidden
- Tip: Use panel() with a tab_group(). A tab_group() is like a popup but is also a container that defaults to place_overlay.
tab_group(bind: @x,
item: [\{name: "tab 1", value: @tab_1\}, \{name: "tab 2", value: @tab_2\}]) \{
panel(bind: @x, value: @tab_1) \{/*...*/ \} panel(bind: @x, value: @tab_2) \{/*...*/ \} \}

\section*{Scripting and Localization}
- Contributing values form the basis for intelligent recording
- Difference between "fixed" values and contributing captures "intent"
- Same model is used for playback - handling all script validation
- Model assists script writers in the same way it assists users - letting them specify the parameters in terms they understand
- ASL contains an experimental xstring library:
button(name: localize("<xstr id='ok'>0K</xstr>"));

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\section*{How Layouts Work}
- A layout is a container of placeable objects
- When a description is parsed a hierarchy of placeable objects is stored in the layout
- The basic algorithm is:
- Gather horizontal metrics of each item in the hierarchy, depth first post order
- Solve the horizontal layout
- Gather vertical metrics - providing final horizontal metrics
- Solve the vertical layout
- Place each item

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\section*{How Layouts Work}


\section*{How Layouts Work}

if \(B\) and \(C\) have compatible guides they collapse to one node

\section*{What you can't do}
- Layouts must be able to be decomposed into rows, columns, and overlays
- No space filling or best fit algorithms

- You can plug-in your own layouts if they can behave as a placeable object.

\section*{How Property Models Work}
- A property model is a container of cells, relationships, views and controllers
- When the description is parsed, cells and relationships are added.
- Views and controller are added from the layout description
- Each cell attached to a relationship has a priority as well as a value, priority is usually based on how recently the element changed

\section*{How Property Models Work}
- The basic algorithm is:
- Calculate the predicates for any conditional relate clauses
- Predicates cannot be involved in relate clauses
- Flow the active relate clauses using the priority on the cells
- After this point, the flow will be use to direct calculations
- Flow and calculate run in opposite directions on the graph.
- Calculate the invariants
- If an invariant is false, any reached source is marked as poison
- Calculate the output expressions
- Reached sources are marked enabled
- If a reached source is poison result is marked invalid
- Calculate any remaining interface cells to which a view is attached

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\section*{What you can't do}
- There are many other types of models that the property model library can't handle - some of the more common ones:
- Sequences (manipulating lists of elements)
- Although the property model can describe invariants on the sequence and preand post- conditions on the functions that manipulate it.
- Grammars
- The property model library is not a parser
- Triggers - imperative actions
" There is no way to say "when this happens do this"
- The property model library cannot handle distributing values (yet)
- From our exercise - there is no way to construct a UI which given a final score calculates how many tds, field goals, and extra points are needed to reach it.

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\section*{Closing Comments}
- Website http://stlab.adobe.com
- Don't be afraid to ask questions - subscribe to our mailing list
- Please contribute to ASL - our charter is to improve how software is written - by contributing you will learn and help others
- We prefer small contributions - contribute the big functions when they become small functions leveraging the rest of the library


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